## Tropical Brill-Noether theory

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## The B(aker)-N(orin) game on graphs



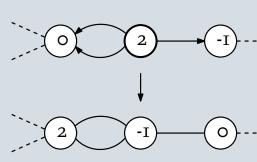


#### Requirements

finite, connected, undirected graph  $\Gamma$   $d \geq 0$  chips natural number r

#### Rules

B puts d chips on  $\Gamma$ N demands  $\geq r_v \geq 0$  chips at v with  $\sum_v r_v = r$ B wins iff he can *fire* to meet N's demand



## Brill-Noether theorems for graphs

$$g:=e(\Gamma)-v(\Gamma)+1 \text{ genus of } \Gamma$$
 
$$\rho:=g-(r+1)(g-d+r)$$

#### **Conjecture (Matthew Baker)**

- 1.  $\rho \ge 0 \Rightarrow B$  has a winning starting position.
- 2.  $\rho < 0 \Rightarrow$  B may not have one, depending on  $\Gamma$ .

 $(\forall g \; \exists \Gamma \; \forall d, r : \rho < 0 \Rightarrow \text{Brill loses.})$ 

#### Theorem (Baker / Caporaso)

I. is true.

(uses sophisticated algebraic geometry)

#### Theorem (Cools-D-Payne-Robeva)

2. is true.

(implies sophisticated algebraic geometry)

## Chip dragging on graphs

Simultaneously moving all chips along edges, with zero net movement around every cycle.

#### Lemma

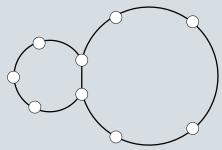
- 1. Chip dragging is realisable by chip firing.
- 2. W.l.o.g. B drags instead of firing.

#### Example 1: $\Gamma$ a tree

$$\begin{array}{l} \rho = g - (r+1)(g-d+r) = -(r+1)(-d+r) \\ \text{B wins} \Leftrightarrow \rho \geq 0 \Leftrightarrow d \geq r \end{array}$$



$$d = 2, r = 1$$
 Who wins?

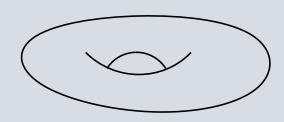


## The B(rill)-N(oether) game on Riemann surface





# Requirements compact Riemann surface X d chips natural number r



#### **Rules**

B puts d chips on XN demands  $\geq r_x \geq 0$  chips at x with  $\sum_x r_x = r$ B wins iff he can drag to meet N's demand

## Chip dragging on Riemann surfaces

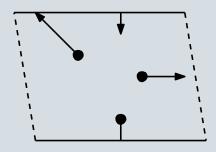
Simultaneously moving chips c along paths  $\gamma_c: [0,1] \to X$ , such that  $\sum_c \langle \omega|_{\gamma(t)}, \gamma'_c(t) \rangle = 0$  for all holomorphic 1-forms  $\omega$  on X.

#### Lemma

$$\begin{split} D &= \sum_c [\gamma_c(0)] \text{ initial position} \\ E &= \sum_c [\gamma_c(1)] \text{ final position} \\ \Leftrightarrow E - D \text{ is divisor of meromorphic function on } X \\ \textit{drag-equivalence} &= \textit{linear equivalence} \end{split}$$

#### **Example: torus**

only one holomorphic 1-form: dz condition:  $\sum_{c} \gamma'_{c}(t) = 0$  when does B win?



#### **Dimension count**

```
\omega_1, \ldots, \omega_q basis of holomorphic 1-forms
 \mathbf{x} = (x_1, \dots, x_d) \in X \times \dots \times X
 v_i \neq 0 tangent vector at x_i
 \rightsquigarrow matrix A_{\mathbf{x}} = (\langle \omega_i, v_i \rangle)_{ij} \in \mathbb{C}^{g \times d}
 (c_1v_1,\ldots,c_dv_d) infinitesimal dragging direction \Rightarrow A(c_1,\ldots,c_d)^T=0
 x winning for B \Rightarrow
 dragging x fills \geq r-dimensional variety
 where \ker A is > r-dimensional
for B to have a winning position, "need" g-d+r d-r(g-d+r) \geq r \Rightarrow \rho = g-(r+1)(g-d+r) \geq r
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### Brill-Noether theorems for Riemann surfaces

Theorem (Meis 1960, Kempf 1971, Kleiman-Laksov 1972)  $\rho \ge 0 \Rightarrow$  B has a winning position.

#### Theorem (Griffiths-Harris 1980)

- 1.  $\rho < 0 \Rightarrow$  B may lose, depending on X. ( $\forall g \exists X \ \forall d, r : \rho < 0 \Rightarrow$  B loses.)
- 2.  $\rho \ge 0$  and X general  $\Rightarrow \rho = \dim\{\text{winning positions modulo dragging}\}$
- 3.  $\rho = 0$  and X general  $\Rightarrow \# = \#$  standard tableaux of shape  $(r+1) \times (g-d+r)$  with entries  $1, 2, \ldots, g$

## **Specialisation**

#### Algebro-geometric (Baker, Caporaso)

dual graph of special fibre applies to arbitrary fields integral starting positions for B

#### Complex-analytic (Mikhalkin-Zharkov)

conceptually simpler? rational starting positions for B  $\{X_t\}_{t\neq 0}$  family of Riemman surfaces  $\Gamma$  for T of ("tropical limit") holomorphic 1-forms on T on T chip dragging on T

#### Theorem

 $D_t$  winning for B on  $X_t$  and  $D_t \to D$  on  $\Gamma$  for  $t \to 0$   $\Rightarrow D$  winning on  $\Gamma$ .

## Consequences of Specialisation

#### Meis/Kempf/Kleiman-Laksov

( $\rho \ge 0$  implies B wins on Riemann surfaces)

 $\Rightarrow$  same statement for  $\Gamma$ .

No combinatorial proof is known!

#### Cools-D-Payne-Robeva

 $(\rho < 0 \Rightarrow B \text{ loses for suitable } \Gamma)$ 

 $\Rightarrow$  same for Riemann surfaces (Griffiths-Harris 1 and 2, and probably 3).

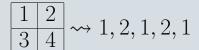
#### **Technical difficulties:**

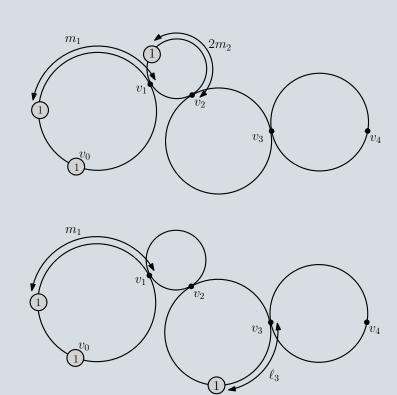
- I. Find family  $\{X_t\}_t$  with dual graph  $\Gamma$  (algebro-geometric) or degenerating to  $\Gamma$  (complex-analytic);
- 2. show that  $t \mapsto D_t$  (winning position on  $X_t$ ) can be chosen such that  $D_t$  "converges".

## Example (Cools-D-Payne-Robeva)

$$g = 4, d = 3, r = 1$$
  
 $g - d + r = 2$   
 $r + 1 = 2, \rho = 0$ 

1	3		1, 2, 3, 2,	1
2	4	7	1, 2, 0, 2,	1

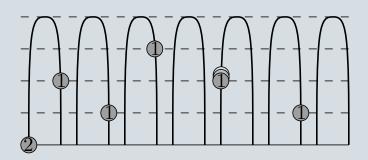




## A larger example

$$g = 7, d = 7, r = 2$$
  
 $\leadsto g - d + r = 2, r + 1 = 3, \rho = 1$ 

1	2	4	000	(21, 31,	, 32, 42, 31,	21 22 21)	lingering lattice	nath
3	6	7	• • • • • • • • • • • • • • • • • • •			$,$ $o_1,$	31, 32, 21)	ingering wince f



#### Theorem (Cools-D-Payne-Robeva)

B's starting position  $\leadsto$  lingering lattice path in  $\mathbb{Z}^r$ ;

B wins iff path stays in chamber  $\{(x_1,\ldots,x_r)\mid x_1>x_2>\ldots>x_r>0\}$ .

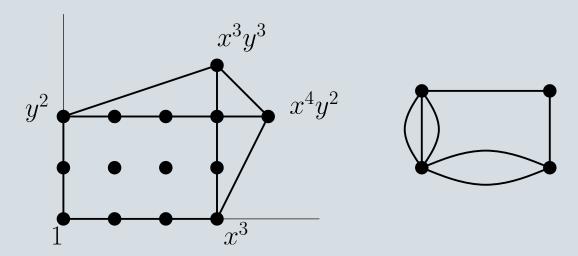
$$\rightsquigarrow \rho \ge 0 \Leftrightarrow B \text{ wins}$$

## Castryck and Cools's gonality conjecture

r=1  $f\in \mathbb{C}[x,y]$  general with Newton polytope  $\Delta$   $X:=\{f=0\}$  Riemann surface

#### Conjecture

minimal d for which B wins on  $X:=\{f=0\}$  (=minimal degree of a meromorphic map to  $\mathbb{P}^1$ ) equals d= lattice width of  $\Delta$  (with two exceptions)

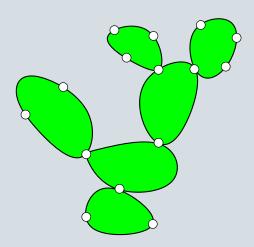


## Purely combinatorial?

#### Theorem (van der Pol)

 $\rho \geq 0$  and  $\Gamma$  a cactus graph

 $\Rightarrow$  B has winning positions with all chips at vertices.

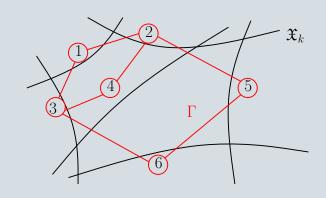


#### Future goal:

Understand Kleiman-Laksov for (metric) graphs.

## Baker's Specialisation Lemma

 $\mathfrak{X}$  curve family over  $\mathbb{C}[[t]]$  (proper, flat, regular scheme) generic fibre  $\mathfrak{X}_{\mathbb{C}((t))}$  smooth curve X special fibre  $\mathfrak{X}_{\mathbb{C}} = X_1 \cup \ldots \cup X_s$   $X_i$  smooth, intersections simple nodes  $\leadsto dual\ graph\ \Gamma$  on  $\{u_1,\ldots,u_s\}$  (metric with edge lengths I)  $\leadsto \max X(\mathbb{C}((t))) \to \{u_1,\ldots,u_s\}$ 



well-behaved with respect to finite extensions  $\mathbb{C}((t^{1/n}))/\mathbb{C}((t))$   $\leadsto$  specialisation map  $\tau: X(\mathbb{C}\{\{t\}\}) \to \Gamma_{\mathbb{Q}}$ 

#### Theorem

Brill wins with starting positing D on  $X(\mathbb{C}\{\{t\}\})$   $\Rightarrow$  Baker wins with starting position  $\tau(D)$  on  $\Gamma_{\mathbb{Q}}$ 

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