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Questions

- What is c(I) for your favourite I?
- Is c(I) computable from generators of I?
- Why should you care?

The curious case of c(I) = 2

Philosophy

[Eisenbud-Sturmfels, 1996]

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For $n \in \mathbb{Z}_{\geq 1}$, $I_n := ((x-z)^2, nx - y - (n-1)z) \subseteq \mathbb{Q}[x, y, z]$ has $c(I_n) = 2$ and $x^n - yz^{n-1}$ is the lowest-degree binomial in I_n .

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 \exists algorithm that, on input $I \subseteq \mathbb{Q}[x_1, \dots, x_n]$, decides $c(I) \stackrel{?}{=} 2$.

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Question: Is c(I) computable?

Special cases: What if n = 1 or if l is linearly generated?

Outline of the Jensen-Kahle-Kathän algorithm:

- Rule out c(I) = 1, pass to $I \subseteq \mathbb{Q}[x_1^{\pm}, \dots, x_n^{\pm}]$.
- If $x^{\alpha} a \cdot x^{0} \in I$, then Trop $(I) \subseteq \alpha^{\perp}$.
- Compute a basis $\alpha_1, \dots, \alpha_m \in \mathbb{Z}^n$ of Trop $(I)^{\perp}$.
- After a monomial coordinate change: $\alpha_i = e_i$, so $x^{\alpha_i} = x_i$.
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- Need to look for binomials in $J := \mathbb{Q}[x_1^{\pm}, \dots, x_m^{\pm}] \cap I$.
- Trop(J) = {0} $\subseteq \mathbb{R}^m$.
- Hence $A := \mathbb{Q}[x_1^{\pm}, \dots, x_m^{\pm}]/J$ is finite-dimensional.
- Let $M_i \in \text{End}(A)$ be multiplication with x_i .
- Then need to find $\alpha \in \mathbb{Z}^m$ s.t. $M_1^{\alpha_1} \cdots M_m^{\alpha_m} = a \cdot id_A$.
- This had already been solved in number theory.

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$$I(X) \subseteq K[x_1, ..., x_n] \leadsto c(X) := c(I(X))$$

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[D-Kahle-Wiersig, 2021]

For a very general r-dimensional *linear space* $X \subseteq K^n$ we have c(X) = r + 1.

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Theorem 2

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For
$$X := \{A \in K^{m \times n} \mid \operatorname{rk}(A) \leq r\}$$
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Theorem 3

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For
$$r$$
 even, $X := \{A \mid A^T = -A, \text{rk}(A) \le r\} \subseteq K^{m(m-1)/2}$ we have $c(X) = (r+1)!! = (r+1)(r-1)(r-3) \cdots 3 \cdot 1$.

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... and in each case we know all $f \in I(X)$ with c(X) terms.

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- Each $f_{\alpha} \neq 0$ vanishes on $\{ rk \leq (r-1) \} \leadsto \geq r!$ terms.

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- Each $f_{\alpha} \neq 0$ vanishes on $\{ rk \leq (r-1) \} \rightsquigarrow \geq r!$ terms.
- If an $f_{\alpha} \neq 0$ vanishes on all rank- $\leq r$ matrices, done.
- Take $A \in K^{(m-1)\times n}$ very general of rank r. Then $f_{\alpha}(A) \neq 0$ for all α with $f_{\alpha} \neq 0$. Theorem 1 applied to the row space X of A yields that $f_{\alpha} \neq 0$ for at least (r+1) distinct α . $\leadsto f$ has at least $(r+1) \cdot r! = (r+1)!$ terms

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Then X defines the uniform matroid of rank 3 on K^6 , but the 4×4 -Pfaffian $x_{12}x_{24} - x_{13}x_{24} + x_{14}x_{23}$ vanishes on X and has 3 < 4 terms.

For a very general *r*-space $X \subseteq K^n$ we have c(X) = r + 1.

• If $F \subseteq K[x_1, ..., x_n]_d$ is a linear subspace spanned by s monomials and some nonzero element of F vanishes on X, then on *each* element of $Gr(r, K^n)$ some nonzero element of F vanishes.

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- Hence $F' \not\subseteq K[x_1, ..., x_r]_d$, as no polynomial in $x_1, ..., x_r$ vanishes on $K^r \times \{0\}^{n-r}$.
- Hence F' contains a monomial x^{β} divisible by some x_j , j > r, but then $(|F| =)|F'| \ge r + 1$.

Characterisation of equality

If $f \in I \subseteq K[x_1, ..., x_n]$, then also:

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Theorem 1'

[D-Kahle-Wiersig, 2021]

For $r \ge 2$, the only (r+1)-term polynomials that vanish on a very general r-space $X \subseteq K^n$ are $c \cdot x^{\alpha} \cdot \ell^{p^e}$ where ℓ is a linear form with r+1 terms.

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Theorem 3'. (r+2)-Pfaffians in the skew-symmetric case.

Key proof ingredient

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Theorem

[Fløystad, 1999]

If $s \ge 3$, char K = 0 and a subspace $F \subseteq K[x_1, ..., x_n]_d$ has $gin_{\succ} F = x_1^{d-1} \cdot \langle x_1, ..., x_s \rangle$, then $F = f \cdot \langle \ell_1, ..., \ell_s \rangle$ for some $f \in K[x_1, ..., x_n]_{d-1}$ and some linear forms $\ell_1, ..., \ell_s$.

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Open problem: the case of symmetric matrices!

For a family of polynomials $S \subseteq K[x_1, ..., x_n]$, a *hitting set* generator is a polynomial map $g: K^m \to K^n$ such that $f \circ g \in K[y_1, ..., y_m] \setminus \{0\}$ for all $f \in S \setminus \{0\}$.

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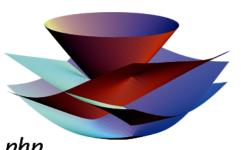
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Observation

[Robert Andrews]

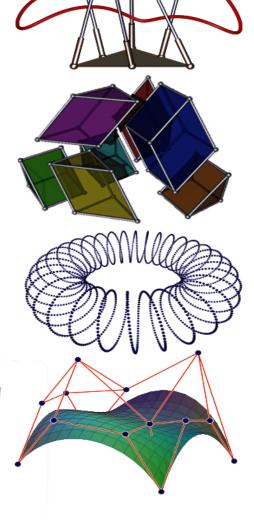
For $S = \{\text{polynomials with} \leq t \text{ terms}\}$, choose r such that $(r+1)! \geq t$, Theorem 2 gives a degree-two hitting set generator $g: K^{\sqrt{n} \times r} \times K^{r \times \sqrt{n}} \to K^{\sqrt{n} \times \sqrt{n}} = K^n$, $(A, B) \mapsto AB$.

The resulting $m = c \cdot \sqrt{n} \cdot \log(t) / \log(\log(t))$ is near optimal.



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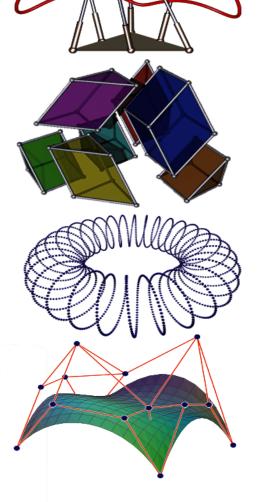








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Thank you!