

A short history of strength

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\rightsquigarrow either $\overline{\mathrm{GL}(U) \times \mathrm{GL}(V)} A = U \otimes V$

or else A is a polynomial in $2k$ *one-tensors*.

Matrix tuple strength=tuple rank

Definition. $A_1, \dots, A_\ell \in \mathbb{C}^{N \times N} \rightsquigarrow \text{rk}(A_1, \dots, A_\ell) := \min\{\text{rk}(c_1 A_1 + \dots + c_\ell A_\ell) \mid (c_1 : \dots : c_\ell) \in \mathbb{P}^{\ell-1}\}$

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Strength of cubics=q-rank

Definition (Derksen-Eggermont-Snowden)

$$f = a_{111}x_1^3 + a_{112}x_1^2x_2 + \cdots + a_{ijk}x_ix_jx_k + \cdots$$

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Theorem (Derksen, Eggermont, Snowden, 2017)

- $\text{q-rank}(f) = \min\{\text{codim}(V) \mid V \subseteq \mathbb{C}^{\oplus\infty}, f|_V = 0\}$
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Then f is a polynomial in $2k$ lower-degree forms.

Strength of higher-degree forms

Definition (Ananyan-Hochster)

$\text{strength}(f = \sum_{i_1 \leq \dots \leq i_d} c_{i_1, \dots, i_d} x_{i_1} \cdots x_{i_d})$ is
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Plan today: generalise the *dense/bounded* dichotomy; applications; computational issues.

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Examples of degree 3:

$$PV = V \otimes V \otimes V, \quad PV = U \oplus V \oplus \wedge^2 V \oplus S^3 V.$$

Polynomial transformations

$$\frac{\text{polynomial maps}}{\mathbb{A}^n} = \frac{?}{\text{polynomial functors}}$$

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Definition. P, Q polynomial functors
 $\alpha : P \rightarrow Q$ *polynomial transformation* if

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Example (rank-one tensors). $PV = V \oplus V \oplus V$,
 $QV = V^{\otimes 3}$, $\alpha_V(v_1, v_2, v_3) = v_1 \otimes v_2 \otimes v_3$.

Polynomial transformations

Definition. $X \subseteq P$ *closed* if $X(V) \subseteq PV$
Zariski-closed and $P(\varphi) : X(U) \rightarrow X(V)$.

Closed subsets form a category; morphisms are restrictions of polynomial transformations.

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Classical

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Proposition (Birk-Eggermont-D-Snowden)

$X \subseteq P$ closed \rightsquigarrow any polynomial transformation
 $X \rightarrow Q$ extends to P .

The dense/bounded dichotomy

Definition (lex order on polynomial functors)

P a poly functor with highest degree part $P_d \supseteq R \supsetneq 0$; $Q < P$ if Q a quotient of $V \mapsto P(U \oplus V)/R$.

Example: $S^2(U \oplus V)/S^2V = S^2U + U \otimes V$.

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Dichotomy (BDES)

(infinite version)

Let $p \in \lim_{\leftarrow n} P(\mathbb{C}^n) =: P_\infty$. Then either $GL_\infty p$ is dense in P_∞ or $\exists Q < P, \alpha : Q \rightarrow P, p \in \alpha(Q_\infty)$.

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Dichotomy (BDES) *(finite version)*

$X \subseteq P$ closed. Then either $X = P$ or $\exists Q_1, \dots, Q_k \prec P$ and $\alpha_i : Q_i \rightarrow P$ such that $X \subseteq \bigcup \text{im}(\alpha_i)$.

(implies all earlier theorems)

First applications

Theorem (D, 2017) (Noetherianity)

Any chain $P \supseteq X_1 \supseteq X_2 \supseteq \dots$ (X_i closed) stabilises.

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Theorem (BDES) (orbit closures)

The map $\alpha \mapsto \overline{\text{im } \alpha}$ is a surjection from $\{\text{polynomial transformations into } P\}$ to $\{\text{closures of } \text{GL}_{\mathbb{N}}\text{-orbits in } P_{\infty}\}$.

The ultimate notion of strength?

Proposition/Definition (BDES)

For $p \in P_\infty$ there is a unique smallest Q such that $p \in \text{im } \alpha$ for some $\alpha : Q \rightarrow P$. Call this Q the *strength* of p .
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Dichotomy

$p \in P_\infty$ has a dense orbit or p has strength $\prec P$.

Stillman's conjecture

Theorem (Erman-Sam-Snowden 2018)

The ring R of bounded-degree series $\sum_{|\alpha| \leq d} c_\alpha x^\alpha$ in variables x_1, x_2, \dots *is a graded polynomial ring* in uncountably many variables of degrees $1, 2, 3, \dots$

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Theorem (Ananyan-Hochster 2016, ESS2017)

There exists an upper bound $N(d_1, \dots, d_k)$ on the projective dimension of (f_1, \dots, f_k) with $\deg(f_i) = d_i$.

A Gröbner proof of Stillman's conjecture

Theorem (D-Lasoń-Leykin, 2018)

Any finitely generated homogeneous ideal in R has a finite grevlex Gröbner basis w.r.t. $x_1 > x_2 > \dots$

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Theorem (D-Lasoń-Leykin, 2018)

There exists an algorithm that on input d_1, \dots, d_k outputs all possible grevlex generic ideals of ideals $(f_1, \dots, f_k) \subseteq R$ where f_i homogeneous of degree d_i .

Final remarks

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